

## CS\_218 COMPILERS

(Attempt 2 questions out of 3)

### Question 1.

- (a) Consider a language which consists of an *assignment* symbol  $:=$ , the four mathematical symbols  $-$ ,  $+$ ,  $*$ ,  $/$ , and *identifiers* formed by an alphabetic character followed by nought or more alphanumeric characters. Write a *Lex* file which recognises and returns **id**, **plus**, **minus**, **mul**, **div**, or **assign** appropriately. (You may assume that the tokens have been declared elsewhere). Your program should contain a **main** function which calls the parser and prints nothing to the screen except the number of lines in the input file.

[8 marks]

- (b) Remove the  $\epsilon$  productions from the following grammar  $G_1$  in such a way that the language of the resulting grammar  $L(G_2) = L(G_1)$ :-

$A \rightarrow Ba|b$   
 $B \rightarrow c|CdD$   
 $C \rightarrow \epsilon|eAD$   
 $D \rightarrow \epsilon|f$

[4 marks]

- (c) Explain the important aspects of the *semantical analysis* phase of an *analysis-synthesis* compiler.

[4 marks]

- (d) Create *first* and *follow* sets for the following grammar :-

$E \rightarrow TE'$   
 $E' \rightarrow +TE'|\epsilon$   
 $T \rightarrow FT'$   
 $T' \rightarrow *FT'|\epsilon$   
 $F \rightarrow (E)|id$

[5 marks]

- (e) State, in words and mathematically, the two disjunctions necessary for a grammar to be LL(1).

[4 marks]

### Question 2.

- (a) Draw a DFA state transition diagram which recognises all strings containing the substring  $ab$  when the alphabet  $\Sigma = \{a, b, c\}$  [5 marks]

- (b) Construct the LR state set for the following augmented grammar :-

$$E' \rightarrow E$$

$$E \rightarrow E + T | T$$

$$T \rightarrow T * F | F$$

$$F \rightarrow (E) | id$$

[8 marks]

- (c) Is the grammar in (b) above LR(0). Give detailed explanations for your answer.

[5 marks]

- (d) Given the following LL(1) parsing table show a trace of the input string  $id + id * id$ .

	id	+	*	(	)	\$
E	$E \rightarrow TE'$			$E \rightarrow TE'$		
$E'$		$E' \rightarrow +TE'$			$E' \rightarrow \epsilon$	$E' \rightarrow \epsilon$
T	$T \rightarrow FT'$			$T \rightarrow FT'$		
$T'$		$T' \rightarrow \epsilon$	$T' \rightarrow *FT'$		$T' \rightarrow \epsilon$	$T' \rightarrow \epsilon$
F	$F \rightarrow id$			$F \rightarrow (E)$		

[7 marks]

### Question 3.

- (a) Remove immediate left recursion from the following grammar :-

$$E \rightarrow E + T | T$$

$$T \rightarrow T * F | F$$

$$F \rightarrow (E) | id$$

[5 marks]

- (b) Given the following LR(0) grammar...

$$S \rightarrow bSe$$

$$S \rightarrow a$$

...and its item state set...

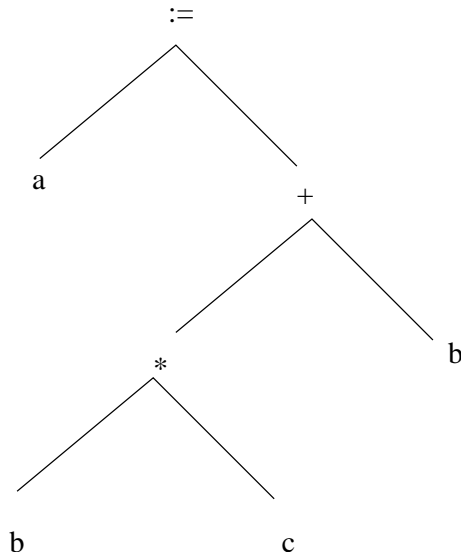
0	$S' \rightarrow \cdot S, S \rightarrow \cdot bSe, S \rightarrow \cdot a$
1	$S' \rightarrow S \cdot$
2	$S \rightarrow b \cdot Se, S \rightarrow \cdot bSe, S \rightarrow \cdot a$
3	$S \rightarrow bS \cdot e$
4	$S \rightarrow a \cdot$
5	$S \rightarrow bSe \cdot$

...construct the LR(0) parse table.

[5 marks]

(c) Let  $G = (T, N, S, P)$  be a *context-free* grammar. Describe each of the four components of  $G$ .  
[4 marks]

(d) Given the following code fragment...  
 $a := b * c + b$   
...and the following tree...



..show what **Three Address Code** might be generated by a compiler. [5 marks]

(e) Draw and carefully label a diagram representing the *analysis-synthesis* model of a compiler. Show on your diagram the parts which are considered as the *front* and *back* ends. Explain why a compiler is usually divided into front and back ends. [6 marks]